



Hyper Parametric Timed CTL

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EMSOFT 2024, 2nd Oct. 2024

Safety Critical CPSs

Study: Single Connected Car Can Trick Smart Traffic Lights Into Causing Intersection Clogging

Waymo robotaxi accident with San Francisco cyclist draws regulatory review

By Reuters

February 9, 2024 5:35 AM GMT+9 · Updated 6 months ago



chigan have shown that on systems can be ehicle. Their research spoofing strategy on one n into believing that an icle-to-infrastructure



on that allows vehicles to share informati utomated toll collection systems, traffic ca

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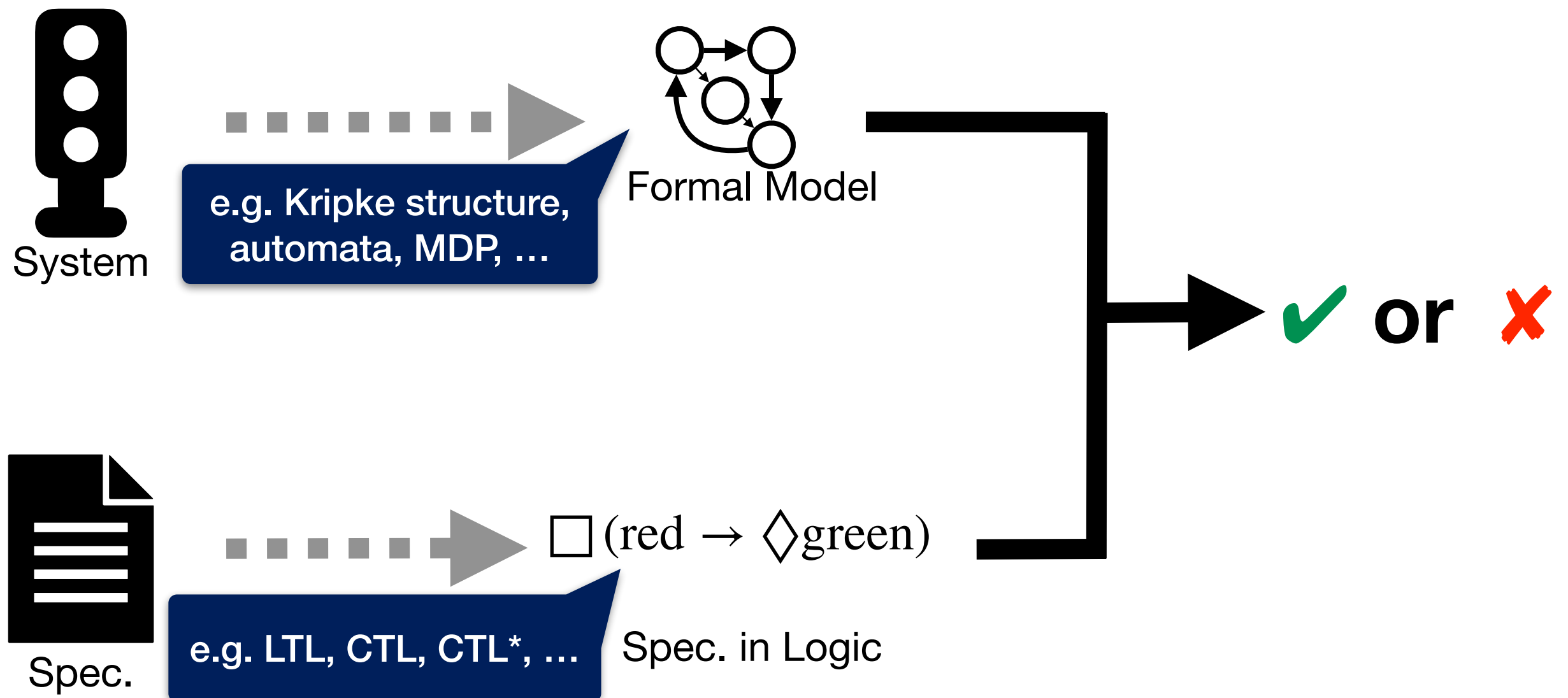
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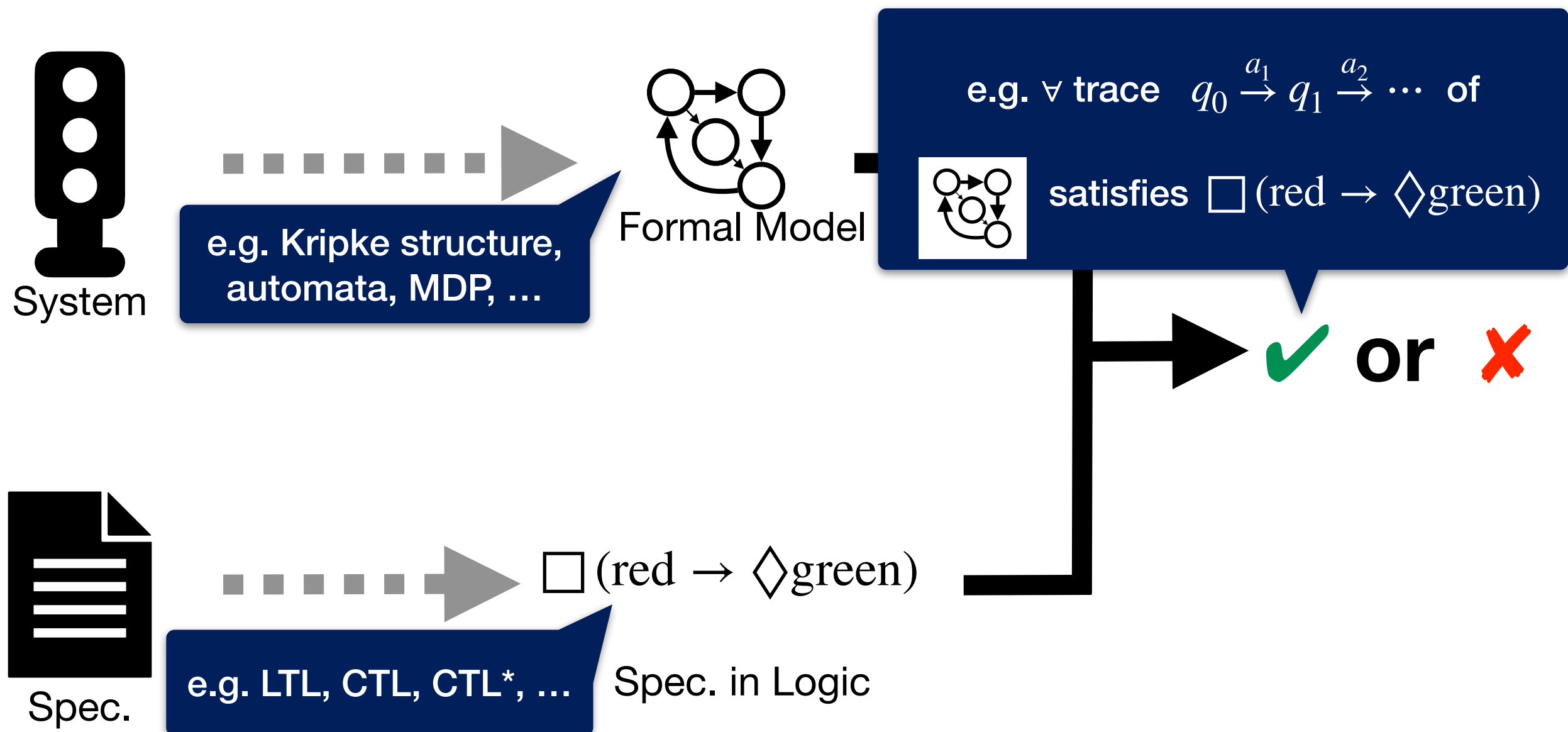
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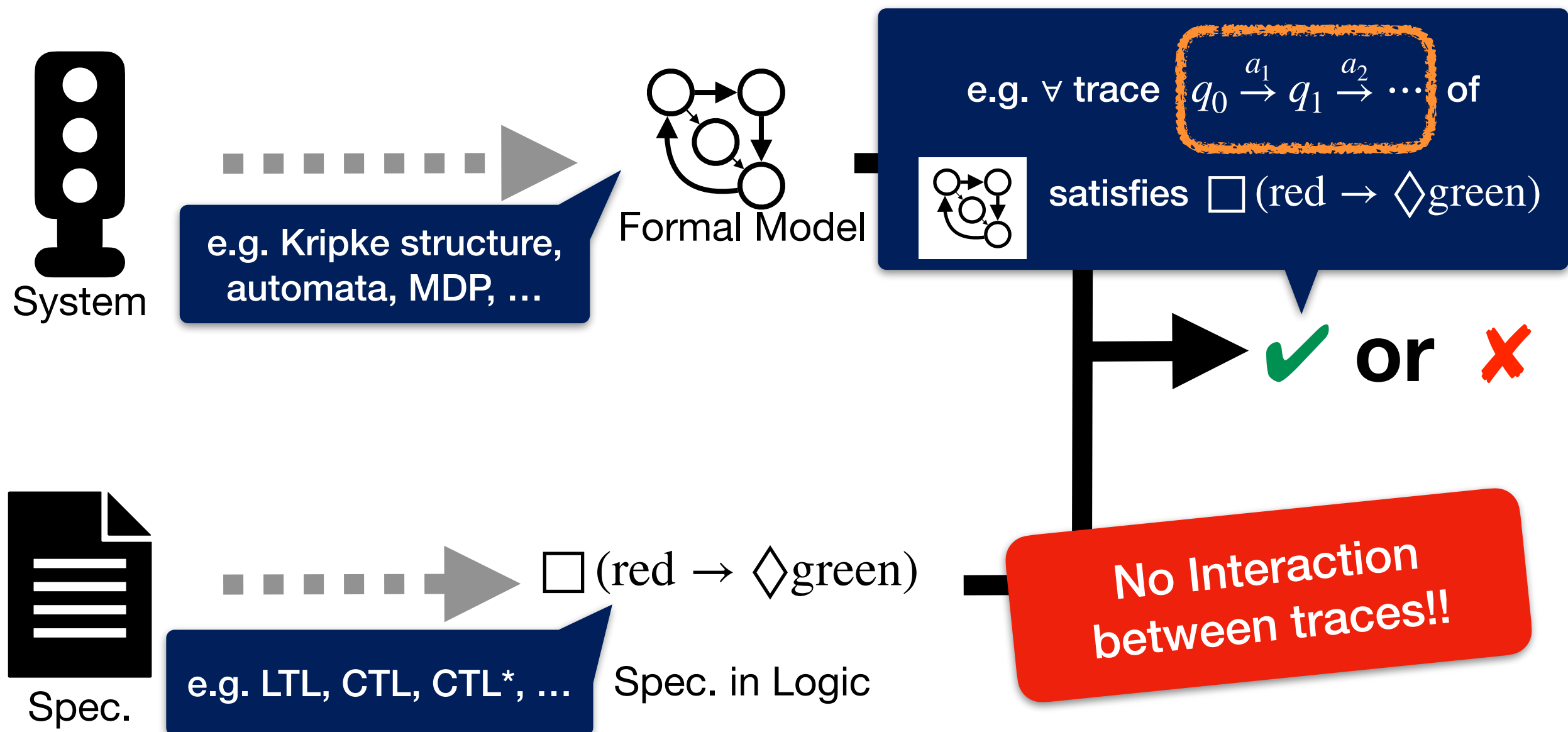
Model Checking



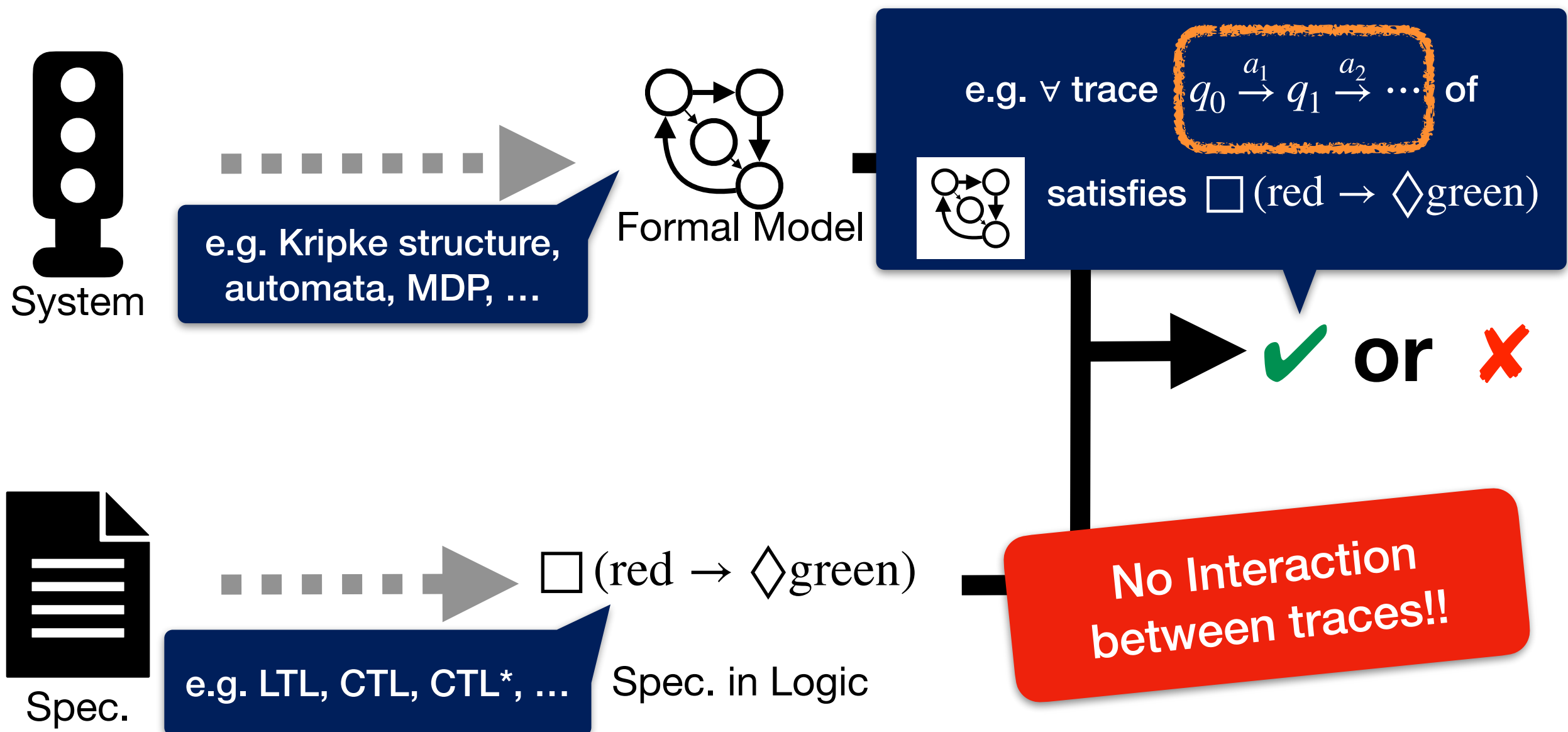
Model Checking



Model Checking



Model Checking of Trace Properties



Hyperproperties

[Clarkson & Schneider, J. Comput. Secur. 2010]

Properties addressing multiple independent traces

Example (Observational determinism)

For the same inputs, the outputs are the same

In other words...

(or $\forall \pi_1, \pi_2 . (\text{Out}_{\pi_1} = \text{Out}_{\pi_2}) \mathcal{W} (\text{In}_{\pi_1} \neq \text{In}_{\pi_2})$ in HyperLTL)

Hyperproperties

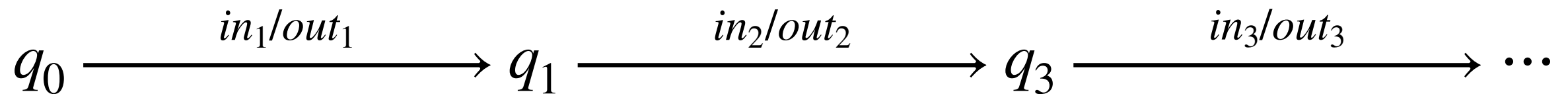
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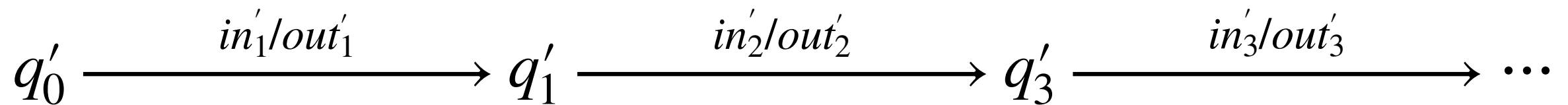
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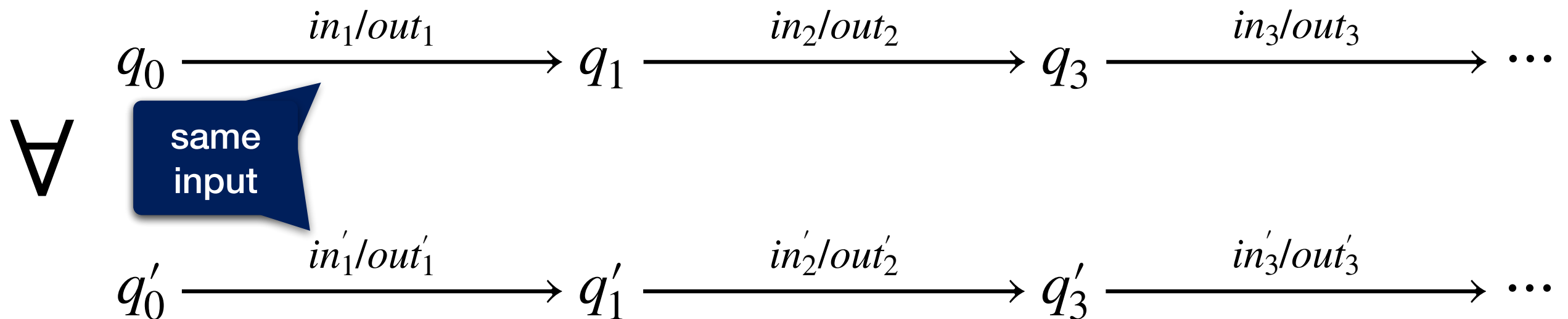
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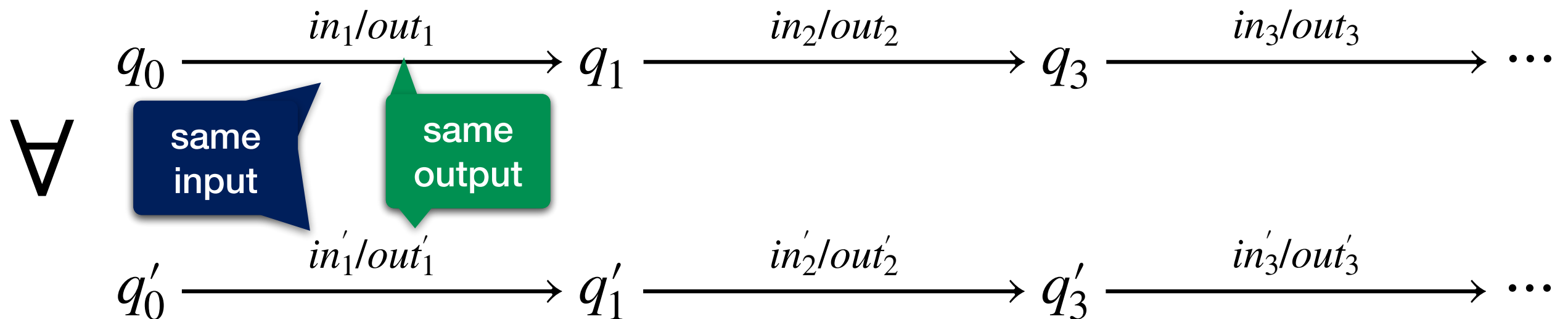
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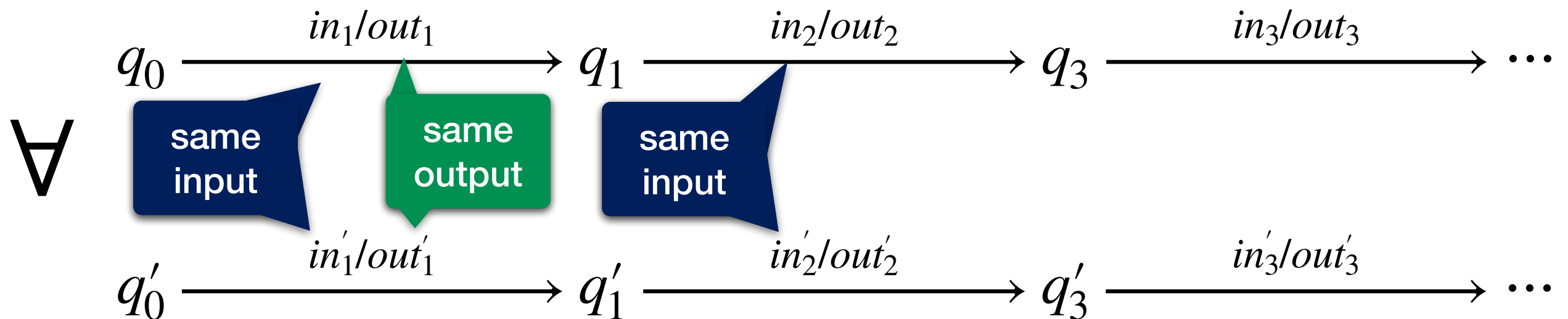
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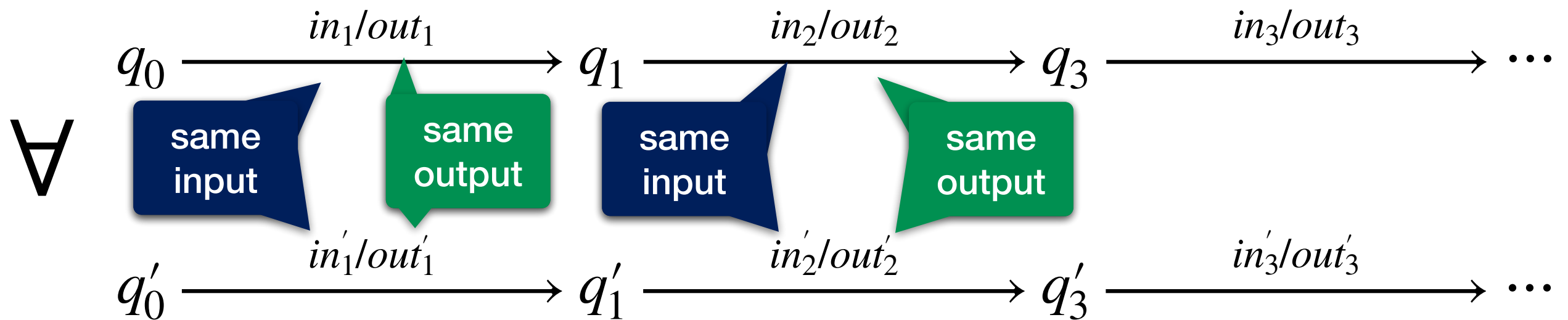
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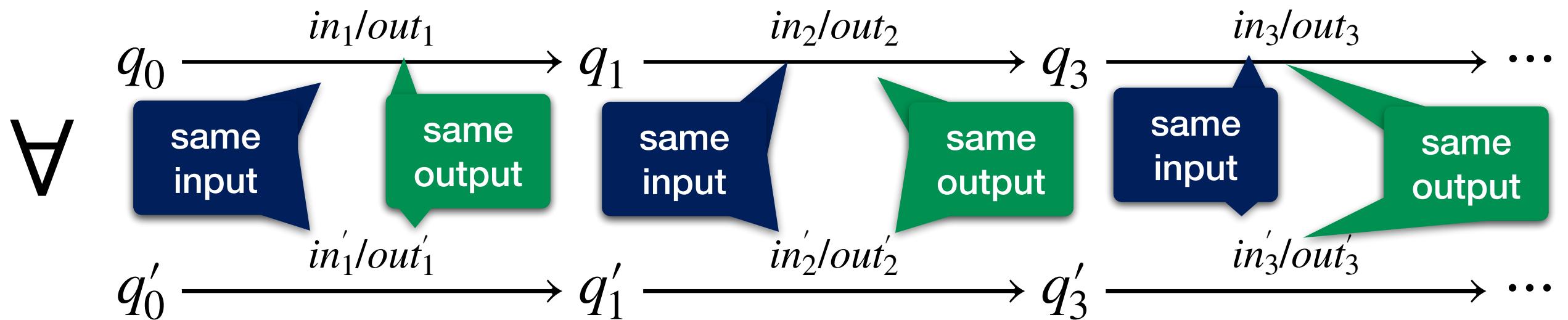
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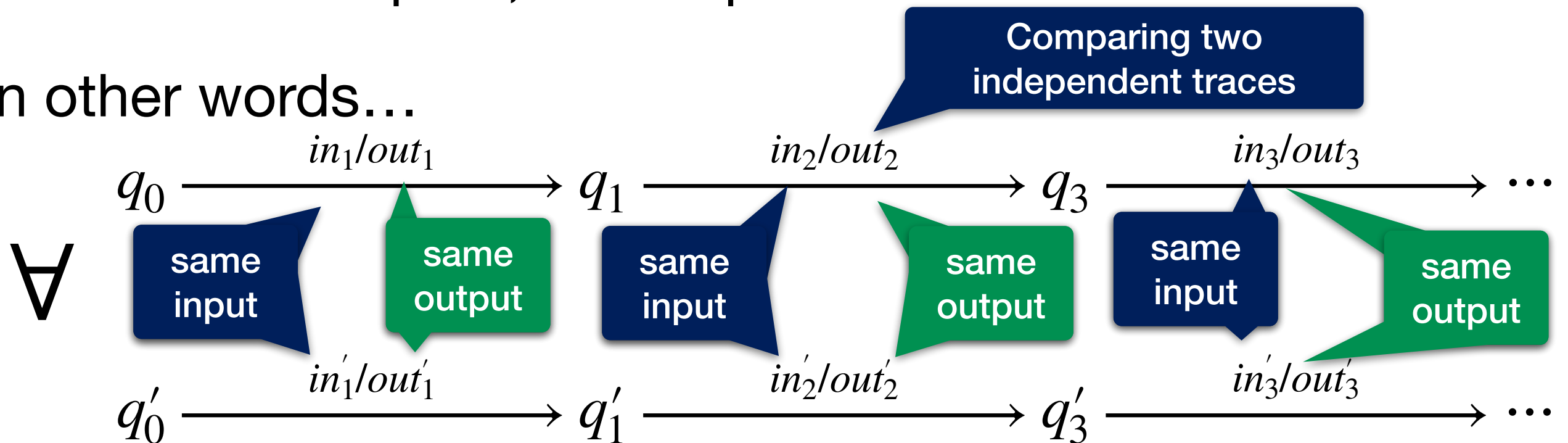
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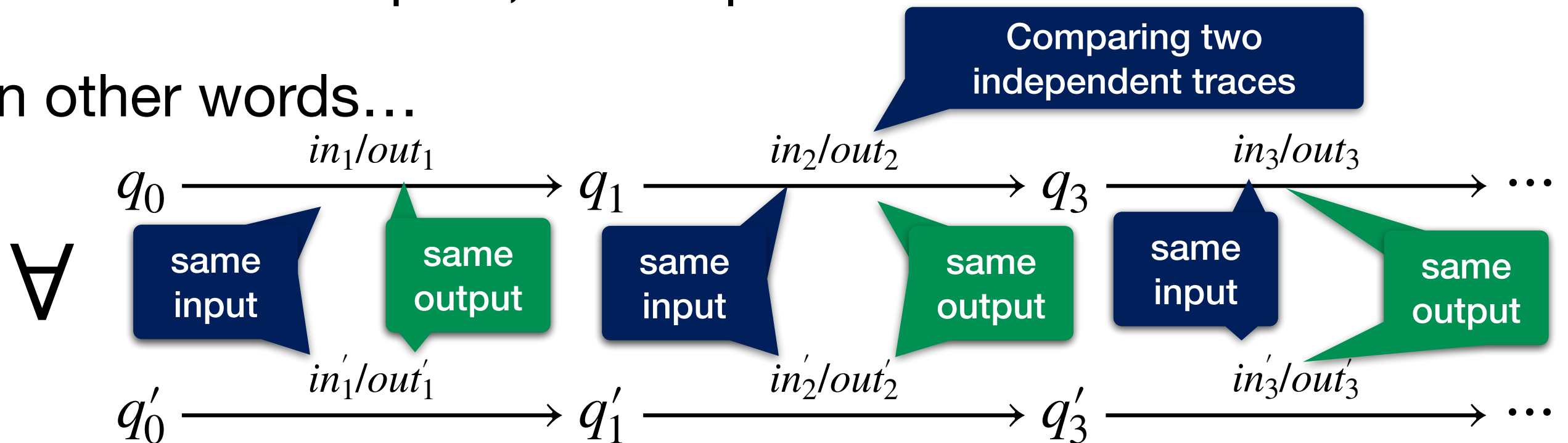
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Useful for Security,
Fairness, Robustness, ...

Q. Can we model check timed hyperproperties?

w/ explicit timing constraints,
continuous notion of time, ...



(Ext-)Hyper Parametric TCTL

[Contribution]

Timed & Parametric Extension of HyperCTL + Additional Predicates

Example (Parametric timed observational determinism)

Observational determinism with small timing deviation of outputs

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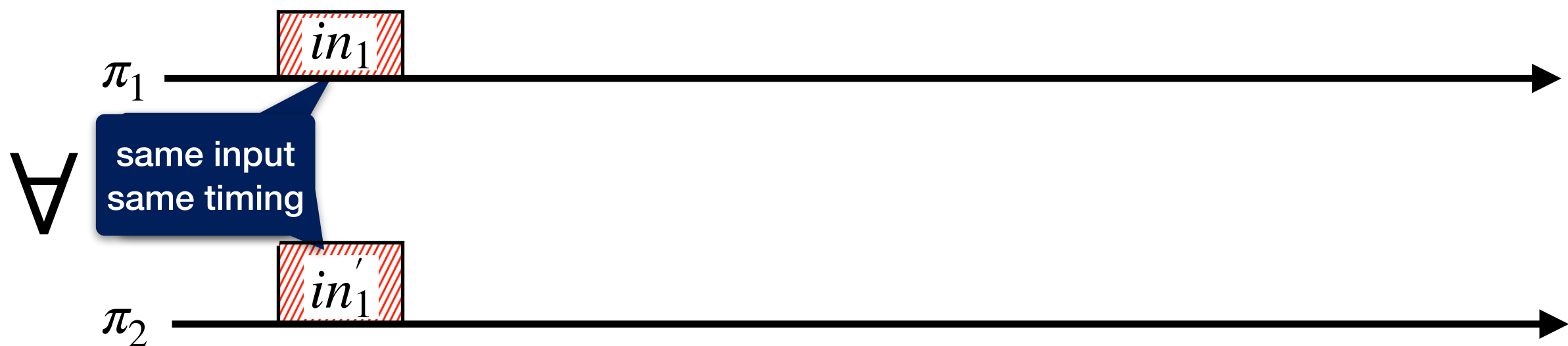
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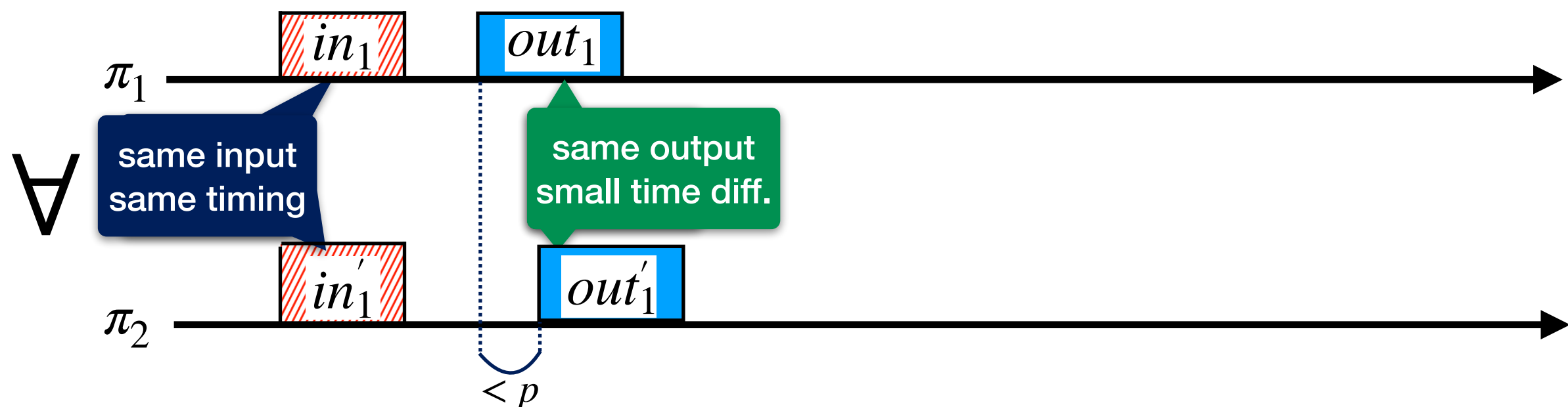
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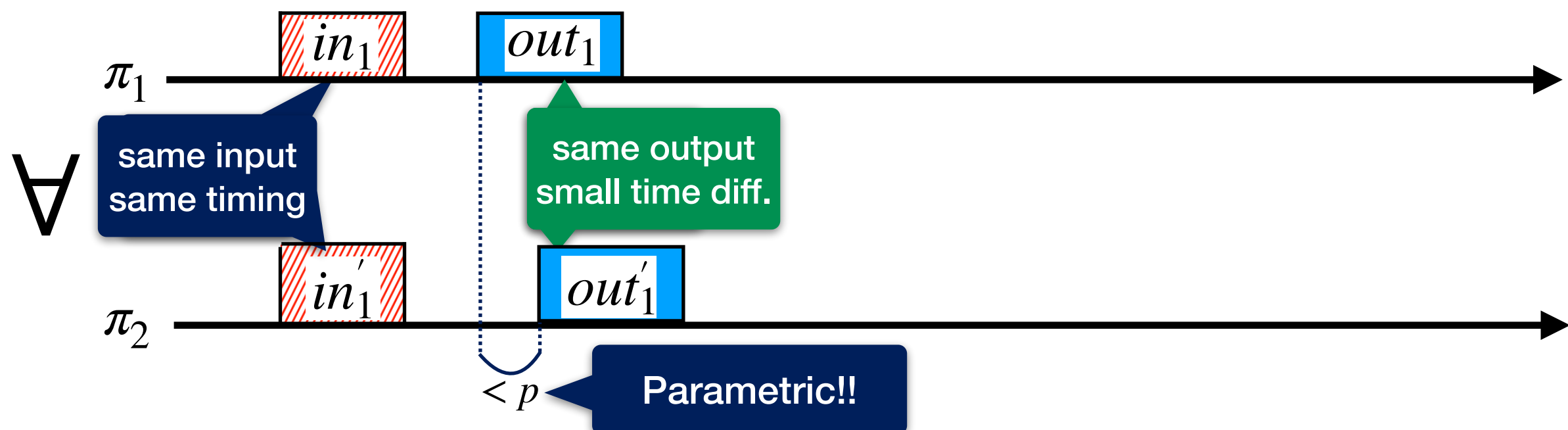
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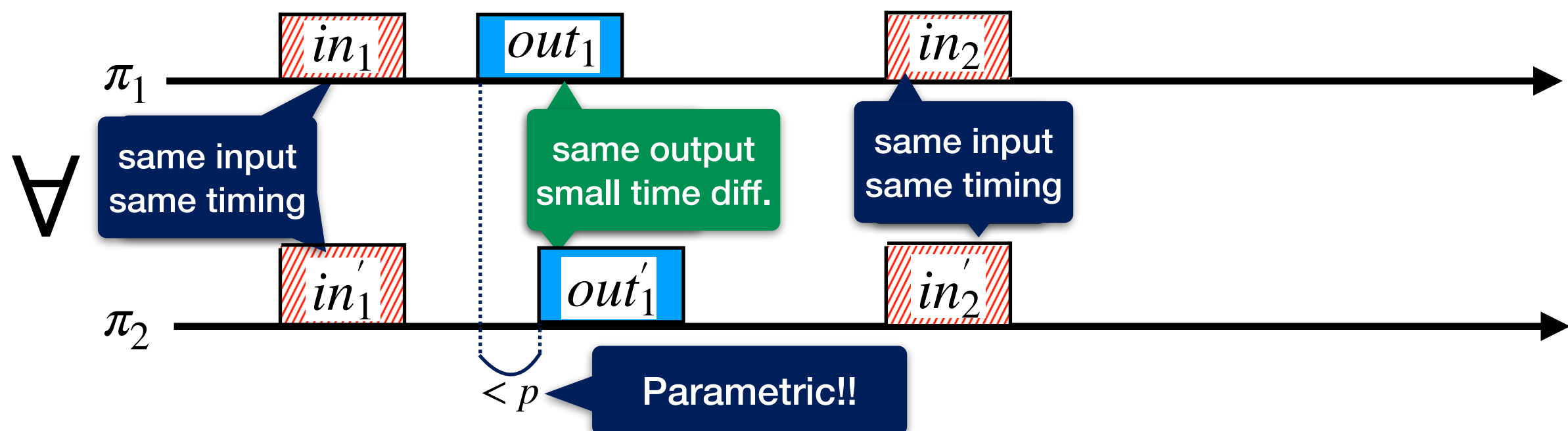
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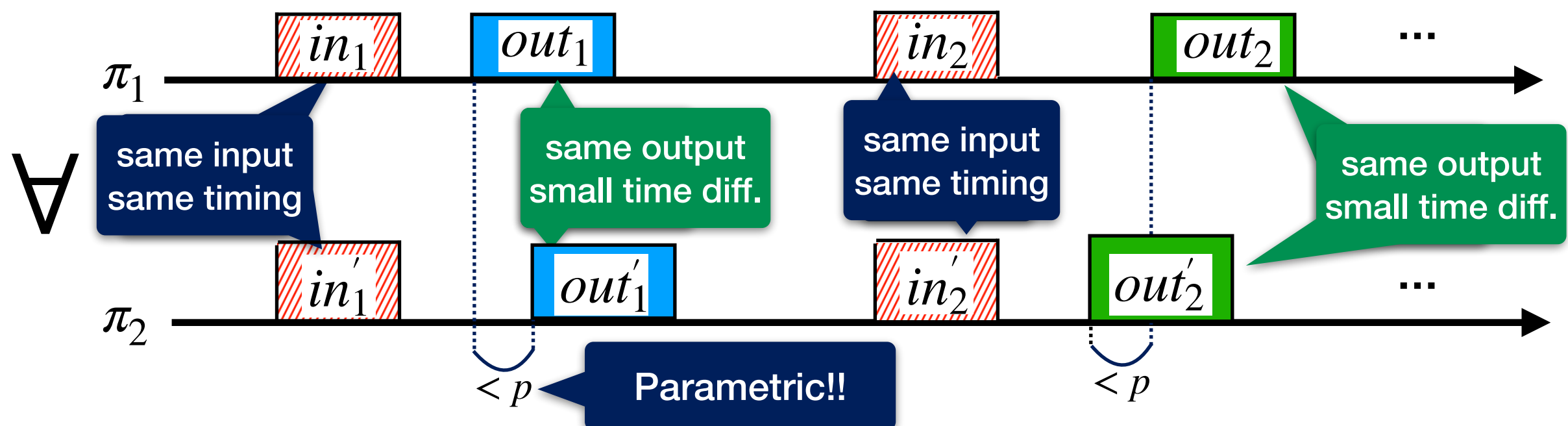
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Ext-Hyper Parametric TCTL

[Henzinger et al., IAC 1994]

Branching + temporal + timing constraints

For any path



is not reachable

$$\varphi = \forall G_{[0,5)} \neg \text{TL}$$

within 5 time units

Ext-Hyper Parametric TCTL

[Bruyère et al., TOCL 2008]

TCTL + parameters in timing constraints

For any path



is not reachable

$$\varphi = \forall G_{[0,p)} \neg \text{TL}$$

within p time units

Ext-Hyper Parametric TCTL

[Contribution]

PTCTL + multiple traces

$$\varphi = \forall \pi_1, \pi_2 (\text{Traffic Light } \pi_1 = \text{Traffic Light } \pi_2) \mathcal{W}_{[0,p)} (\text{Car } \pi_1 \neq \text{Car } \pi_2)$$

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PTCTL + multiple traces

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Timing of  is same in π_1 and π_2

Ext-Hyper Parametric TCTL

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[Contribution]

HyperPTCTL + additional predicates (# and *Last*)

$$\varphi = \forall \pi_1, \pi_2 (\# \text{ 🚦 }_{\pi_1} = \# \text{ 🚦 }_{\pi_2} \Rightarrow |Last(\text{ 🚦 }_{\pi_1}) - Last(\text{ 🚦 }_{\pi_2})| < p)$$

$\mathcal{W}_{[0,p)}(\text{ 🚗 }_{\pi_1} \neq \text{ 🚗 }_{\pi_2})$

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Timing of last  is similar in π_1 and π_2

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$\mathcal{W}_{[0,p)}(\text{Car} \pi_1 \neq \text{Car} \pi_2)$

Timing of  is different in π_1 and π_2

Ext-HyperPTCTL Model Checking

[Contribution]

Idea: Reduction to PTCTL model checking

1. Ext-HyperPTCTL \rightarrow HyperPTCTL: encode w/ PTAs
2. HyperPTCTL \rightarrow PTCTL: self-composition of PTAs

Note: PTCTL model checking is in general undecidable
 \rightarrow We only have semi-algorithm

Ext-HyperPTCTL \rightarrow HyperPTCTL

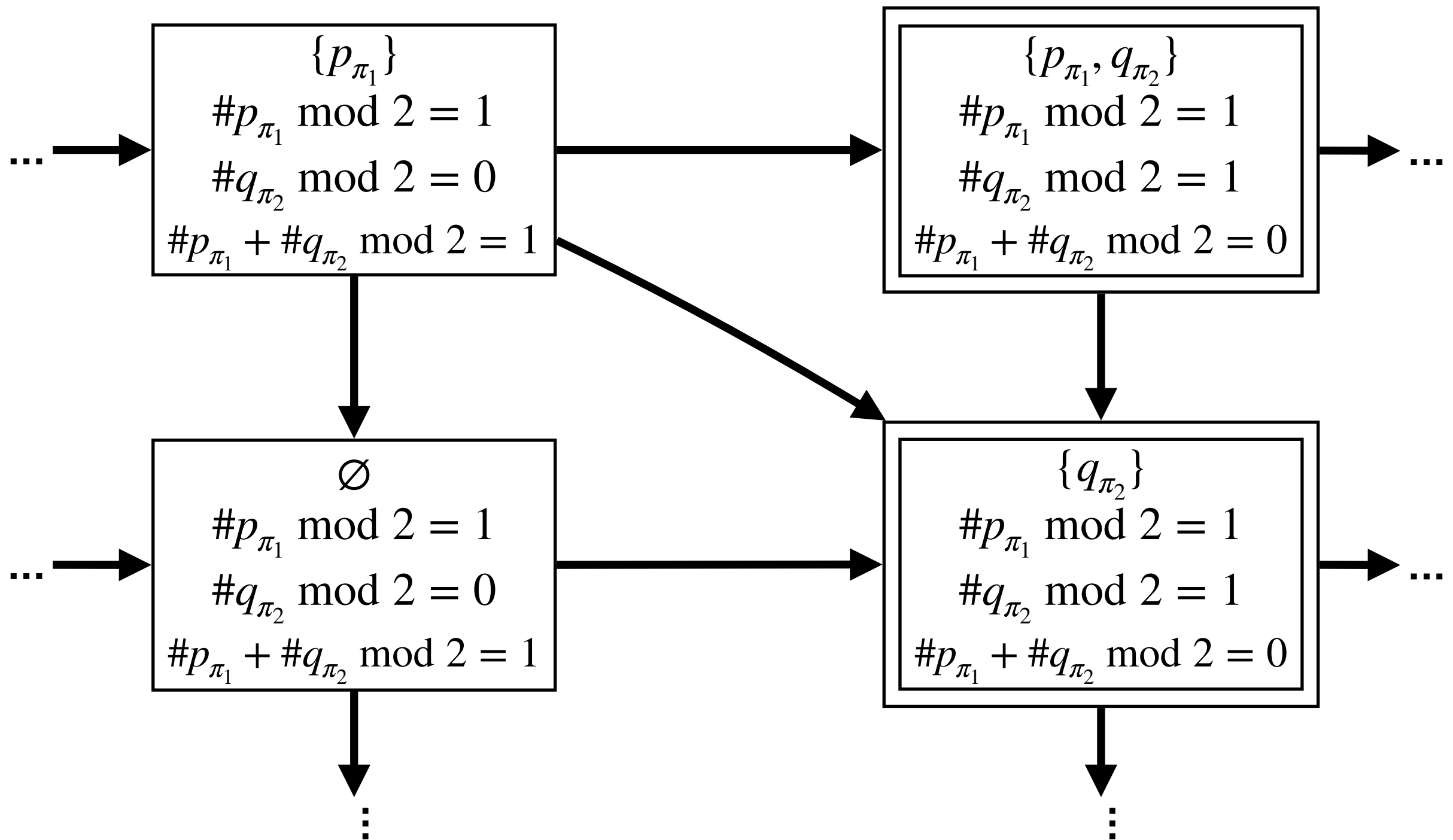
Encode extended predicates w/ observer PTAs

Encoding of general terms with PTAs is hard, but possible for some forms of terms

Slogan: Restrict the terms so that:

- Predicates' truth values are updated only at transitions
- Only finite (discrete) counting is sufficient

Example: $\#p_{\pi_1} + \#q_{\pi_2} \pmod 2 = 0$

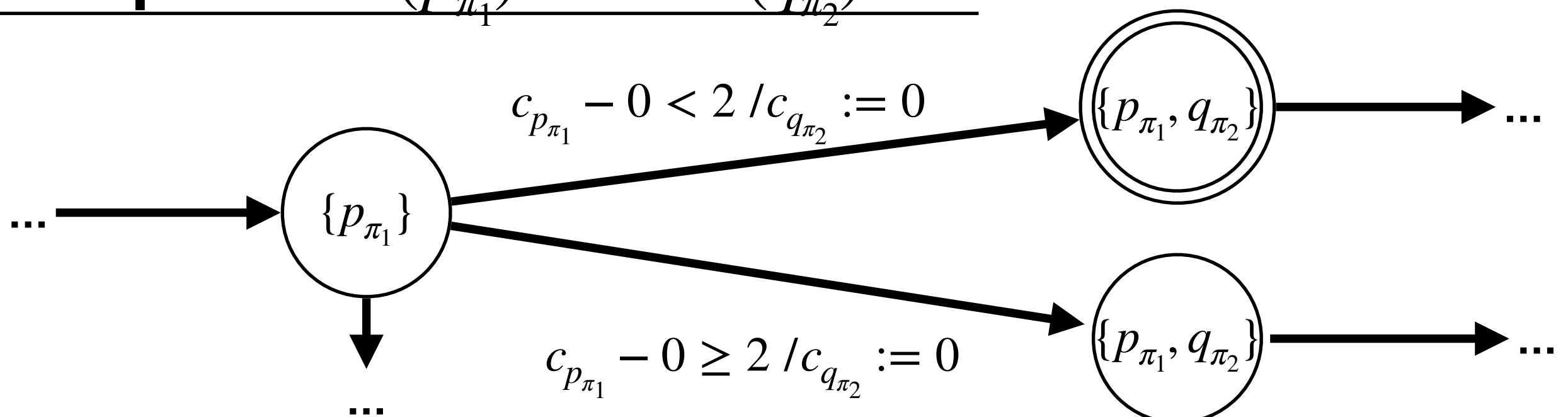


Terms with *LAST*

Linear term over param.

- We support $LAST(p_{\pi_1}) - LAST(p'_{\pi_2}) \bowtie \mu$
→ Truth values changes only at transitions
- Encode $LAST(p_{\pi_1})$ and $LAST(p'_{\pi_2})$ with clocks

Example: $LAST(p_{\pi_1}) - LAST(q_{\pi_2}) < 2$



HyperPTCTL \rightarrow PTCTL

Self-composition to enc. multiple traces by one trace

- Reduction s.t.
traces $\pi_1, \pi_2, \dots, \pi_n$ of \mathcal{A} is captured by
trace $\pi_1 || \pi_2 || \dots || \pi_n$ of $\mathcal{A} || \mathcal{A} || \dots || \mathcal{A}$

Limitation/Challenge

Quantifier alternation \approx autom. complement

- Complement is impossible
 \rightarrow Focus on nest-free fragment
- “zero-time” behavior is tricky

Limited, but still likely useful

Challenge: Zero-time behavior of TA

Multiple transitions may occur at the same time

$$\pi_1 = (l_0, \nu_0) \xrightarrow{\text{jump}_1} (l_1, \nu_1) \xrightarrow{\text{jump}_2} (l_2, \nu_2)$$

$$\pi_2 = (l'_0, \nu'_0) \xrightarrow{\text{jump}'_1} (l'_1, \nu'_1) \xrightarrow{\tau=2.4} (l'_2, \nu'_2)$$

Jumps in $\pi_1 \parallel \pi_2$ can be ...

- $\text{jump}_1 \rightarrow \text{jump}_2 \rightarrow \text{jump}'_1$
- $\text{jump}_1 \rightarrow \text{jump}'_1 \rightarrow \text{jump}_2$
- $\text{jump}_1 \parallel \text{jump}'_1 \rightarrow \text{jump}_2$
- ...

Explicit Transition Ordering

Idea: Path valuation := (paths, order)

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Transition ordering: $\text{jump}_1 \sim \text{jump}'_1 < \text{jump}_2$

Jumps in $\pi_1 \parallel \pi_2$ is

- $\text{jump}_1 \parallel \text{jump}'_1 \rightarrow \text{jump}_2$

Implementation/Experiments

Model checker for PTAs

- Implemented the reduction to IMITATOR
 - HyPTCTLchecker: Published under GPLv3
- Reduction slightly differs from theoretical one
e.g. IMITATOR's discrete var. not encoding w/ loc.
- The reduction is almost immediate
→ Report the result of synthesis with IMITATOR
- AWS EC2 m7i.4xlarge (16vCPU and 64 GiB RAM) w/ Ubuntu 22.04 LTS.

Experimental Results

Prop. (ψ)	PTA (\mathcal{A})	$ L $	$ C $	$ \mathbb{P} _\psi$	$ \mathbb{P} _{\mathcal{A}}$	$ \mathcal{V} $	Time [sec.]
Deviation	ClkGen	2	1	1	1	2	4.116
Opacity	Coffee	6	2	0	3	2	0.723
Opacity	STAC1:n	8	2	0	2	2	0.178
Opacity	STAC4:n	9	2	0	5	2	< 0.001
Unfair	FIFO	63	2	0	4	2	71.955
Unfair	Priority	72	2	0	4	2	6.855
Unfair	R.R.	81	3	0	4	2	12550.979
RobOND	Coffee	6	2	1	3	2	3.182
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Parametric Unfairness
for Round-Robin
Scheduler < 3.5 h

Slow for large models

Experimental Results

Prop. (ψ)	PTA (\mathcal{A})	$ L $	$ C $	$ \mathbb{P} _\psi$	$ \mathbb{P} _{\mathcal{A}}$	$ \mathcal{V} $	Time [sec.]
Deviation	ClkGen	2	1	1	1	2	4.116
Opacity	Coffee	6	2	0	3	2	0.723
Opacity	STAC1:n	8	2	0	2	2	0.178
Opacity	STAC4:n	9	2	0	5	2	< 0.001
Unfair	FIFO	63	2	0	4	2	71.955
Unfair	Priority	72	2	0	4	2	6.855
Unfair	R.R.	81	3	0	4	2	12550.979
RobOND	Coffee	6	2	1	3	2	3.182
RobOND	WFAS ₀ ¹	24	4	1	0	2	1.665
RobOND	WFAS ₀ ²	24	4	1	0	2	2.570
RobOND	WFAS ₁	24	4	1	1	2	67.644
RobOND	WFAS ₂	24	4	1	2	2	1332.310
RobOND	ATM	7	2	1	0	2	T.O.
RobOND	ATM'	5	2	1	0	2	4179.197
EF ₂	Coffee	6	2	1	0	2	0.034
EF ₃	Coffee	6	2	1	0	3	159.541
EF ₄	Coffee	6	2	1	0	4	T.O.

Scalability wrt # of path
is not good
due to self-composition

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Can handle 3 paths

Related Logics/Methods

	Logic	Continuous Time?	Parametric?	Model Checking?	Quantifier Alternation in Model Checking?
[Ours]	Ext-HyperPTCTL	Yes	Yes	Yes	No
[Ho+, 2021]	HyperMITL	Yes	No	Yes	Only for untimed models
[Bonakdarpour+, 2020]	HyperMTL	No	No	Yes	Yes
[Bartocci+, 2023]	HyperSTL	Yes	Yes	No (Req. Mining)	N/A

Conclusions

- Introduce HyperPTCTL and Ext-HyperPTCTL
- Semi-algorithm for model checking/synthesis of PTAs against nest-free Ext-HyperPTCTL
- Show decidable subclasses even w/ param.
- Implementation & Experiments
 - Works for mild size PTAs

Appendix

Syntax of (Ext-)HyperPTCTL

HyperPTCTL

proposition on loc.
(Edge labels are simplification for illustration)

$$\varphi ::= \top \mid \sigma_\pi \mid \neg\varphi \mid \varphi \vee \varphi \mid \exists \pi_1, \pi_2, \dots, \pi_n. \varphi \mathcal{U}_{\Delta \gamma} \varphi$$

$$\mid \forall \pi_1, \pi_2, \dots, \pi_n. \varphi \mathcal{U}_{\Delta \gamma} \varphi$$

temporal level

$$\psi ::= \varphi \mid p \Delta lt_{\geq 0} \mid \neg\psi \mid \psi \vee \psi \mid \tilde{\exists} p \psi$$

top level

Ext-HyperPTCTL

$$\varphi ::= \top \mid \sigma_\pi \mid LAST(\sigma_\pi) - LAST(\sigma_\pi) \Delta lt \mid cnt_{\geq 0} \Delta d$$

$$\mid (cnt \bmod N) \Delta d \mid \dots$$

Semantics of HyperPTCTL

- $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \sigma_{\pi}$ if $\pi \in \mathbf{dom}(\Pi)$ and $\sigma \in \Lambda(\mathit{Init}(\Pi(\pi)))$;
- $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \neg\varphi$ if we have $(\Pi, \trianglelefteq_{\Pi}), s \not\models_{v, \mathcal{A}} \varphi$;
- $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \varphi_1 \vee \varphi_2$ if $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \varphi_1$ or $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \varphi_2$ holds;
- $(\Pi, \trianglelefteq_{\Pi}), s \models_{v, \mathcal{A}} \exists\pi_1, \pi_2, \dots, \pi_n. \varphi_1 \mathcal{U}_{\bowtie\gamma} \varphi_2$ if for some extension $(\Pi^1, \trianglelefteq_{\Pi^1})$ of $(\Pi, \trianglelefteq_{\Pi})$ satisfying $\mathbf{dom}(\Pi^1) = \mathbf{dom}(\Pi) \sqcup \{\pi_1, \pi_2, \dots, \pi_n\}$ and $\Pi^1(\pi_i) \in \mathit{Paths}(v(\mathcal{A}), s)$ for each $i \in \{1, 2, \dots, n\}$, there is $(\Pi^2, \trianglelefteq_{\Pi^2})$ satisfying $(\Pi^1, \trianglelefteq_{\Pi^1}) \succeq (\Pi^2, \trianglelefteq_{\Pi^2})$, $\mathit{Dur}(\Pi^1 - \Pi^2) \bowtie v(\gamma)$, $(\Pi^2, \trianglelefteq_{\Pi^2}), \mathit{Init}(\Pi^2(\pi_n)) \models_{v, \mathcal{A}} \varphi_2$, and for any $(\Pi^3, \trianglelefteq_{\Pi^3})$ satisfying $(\Pi^1, \trianglelefteq_{\Pi^1}) \succeq (\Pi^3, \trianglelefteq_{\Pi^3}) \succ (\Pi^2, \trianglelefteq_{\Pi^2})$, $(\Pi^3, \trianglelefteq_{\Pi^3}), \mathit{Init}(\Pi^3(\pi_n)) \models_{v, \mathcal{A}} \varphi_1$ holds.

$\forall \mathcal{U}$ is omitted

Terms with

- Encode term's values with discrete states
- $\#p_{\pi_1}$ changes only when satisfied prop. changes on π_1
→ Truth values changes only at transitions
- $((\text{linear term with } \#) \bmod N) \bowtie d$: Encode LHS values with N states for each $\#p_{\pi}$
- $(\text{positive linear with } \#) \bowtie d$: LHS strictly increases
→ merge $> d$

Like the abstraction for regions

Terms with *LAST*

Linear term over param.

- We support $LAST(p_{\pi_1}) - LAST(p'_{\pi_2}) \bowtie \mu$
→ Truth values changes only at transitions
- Encode $LAST(p_{\pi_1})$ and $LAST(p'_{\pi_2})$ with clocks

Specifications in Experiments (1/3)

Deviation:

$\exists \pi_1, \pi_2 . (\text{AtMostOneDiff}) \mathcal{U}_{\geq 0} (\text{SameCount} \wedge \text{LargeDeviation})$

- AtMostOneDiff: $(\#H_{\pi_1} - \#H_{\pi_2}) \bmod 4 \in \{0, 1, 3\}$
- SameCount: $(\#H_{\pi_1} - \#H_{\pi_2}) \bmod 4 = 0$
- LargeDeviation: $\text{Last}(H_{\pi_1}) - \text{Last}(H_{\pi_2}) \notin [-p, p]$

Opacity:

$\exists \pi_1, \pi_2 . (\neg \text{Goal}_{\pi_1} \wedge \neg \text{Goal}_{\pi_2}) \mathcal{U}_{=p} (\text{Goal}_{\pi_1} \wedge \text{Goal}_{\pi_2} \wedge \#Private_{\pi_1} = 0 \wedge \#Private_{\pi_2} > 0)$

Specifications in Experiments (2/3)

Unfair: $\exists \pi_1, \pi_2 . (\text{SyncSub}) \mathcal{U}_{\geq 0}(\text{SameCount} \wedge \text{LargeDeviation})$

- SyncSub: $\text{Sub}_{\pi_1}^1 = \text{Sub}_{\pi_2}^2$
- SameCount: $(\#\text{Run}_{\pi_1}^1 - \#\text{Run}_{\pi_2}^2) \bmod 4 = 0$
- LargeDeviation: $\text{Last}(\text{Run}_{\pi_1}^1) - \text{Last}(\text{Run}_{\pi_2}^2) \notin (-5, 5)$

RobOND: $\exists \pi_1, \pi_2 . (\text{SyncIn} \wedge \text{AtMostOneDiff}) \mathcal{U}_{\geq 0}(\text{LargeDeviation})$

- SyncIn: $\forall i \in \{1, 2, \dots, m\} . \text{In}_{\pi_1}^i = \text{In}_{\pi_2}^i$
- AtMostOneDiff: $\forall j \in \{1, 2, \dots, n\} . (\#\text{Out}_{\pi_1}^j - \#\text{Out}_{\pi_2}^j) \bmod 4 \in \{0, 1, 3\}$
- LargeDeviation:
 $\exists j \in \{1, 2, \dots, n\} . (\#\text{Out}_{\pi_1}^j - \#\text{Out}_{\pi_2}^j) \bmod 4 = 0 \wedge \text{Last}(\text{Out}_{\pi_1}^j) - \text{Last}(\text{Out}_{\pi_2}^j) \notin [-p, p]$

Specifications in Experiments (3/3)

Not supported in theory, but supported in our implementation w/ IMITATOR's discrete variables

$$\underline{\text{EF}}_i: \exists \pi_1, \pi_2, \dots, \pi_i \cdot \diamond_{[p,p]} \forall j \in \{1, 2, \dots, i\} \#a_{\pi_j} - \#a_{\pi_{j+1}} = 1$$